



Portland State University

Winter'21 CS 584/684

Algorithm Design & Analysis

What is an **algorithm**?



In **mathematics** and **computer science**, an **algorithm** ([/ˈælgərɪðəm/](#) [\(listen\)](#)) is a finite sequence of **well-defined**, computer-implementable instructions, typically to solve a class of problems or to perform a computation.^{[\[1\]](#)[\[2\]](#)} Algorithms are always **unambiguous** and are used as specifications for performing **calculations**, **data processing**, **automated reasoning**, and other tasks.

Can you name a few algorithms?

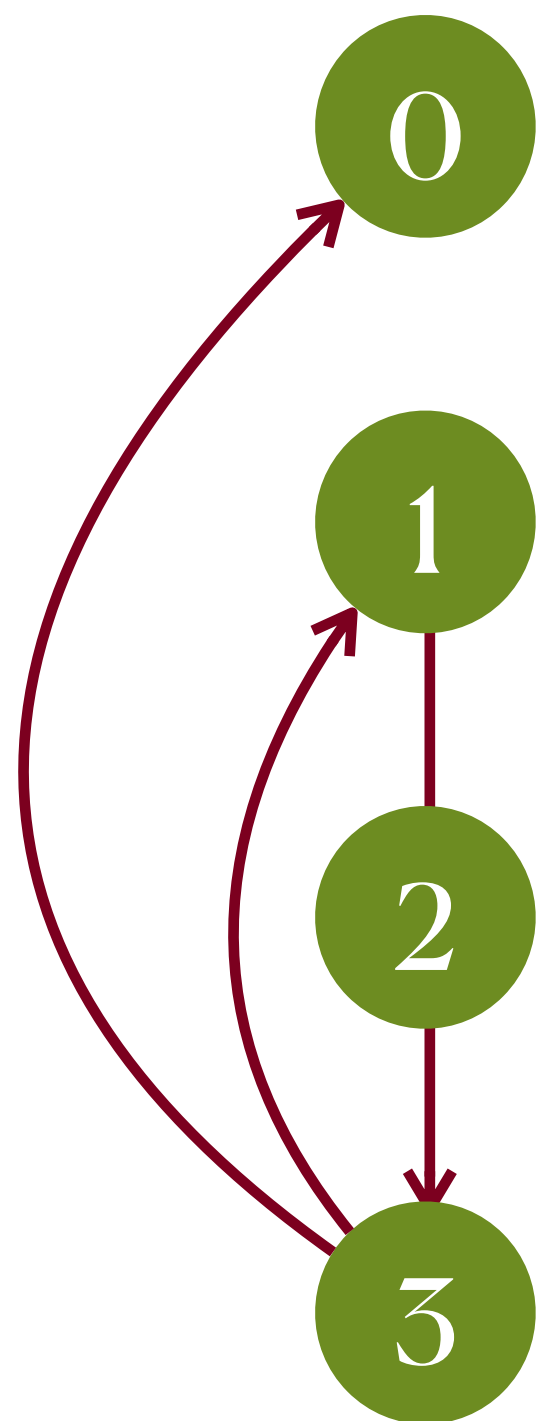
Multiplication: back to grade school

© Long multiplication algorithm

$$\begin{array}{r} 2021 \\ \times 365 \\ \hline 10105 (2021 \times 5) \\ 121260 (2021 \times 60) \\ + 606300 (2021 \times 300) \\ \hline 737665 \end{array}$$

- Unambiguous set of instructions: **above**
- Solving a well-defined problem: **multiplying two non-negative integers**

Principal questions



What **problem** to solve?

Is the algorithm **correct**?

How much **resource** it costs?

Can we do **better**?

DESIGN

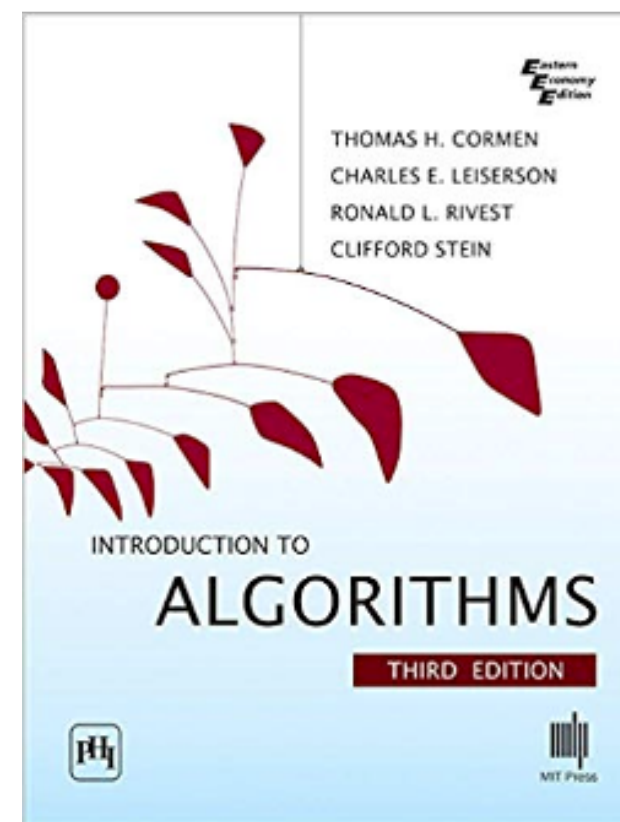
ANALYSIS

Logistics

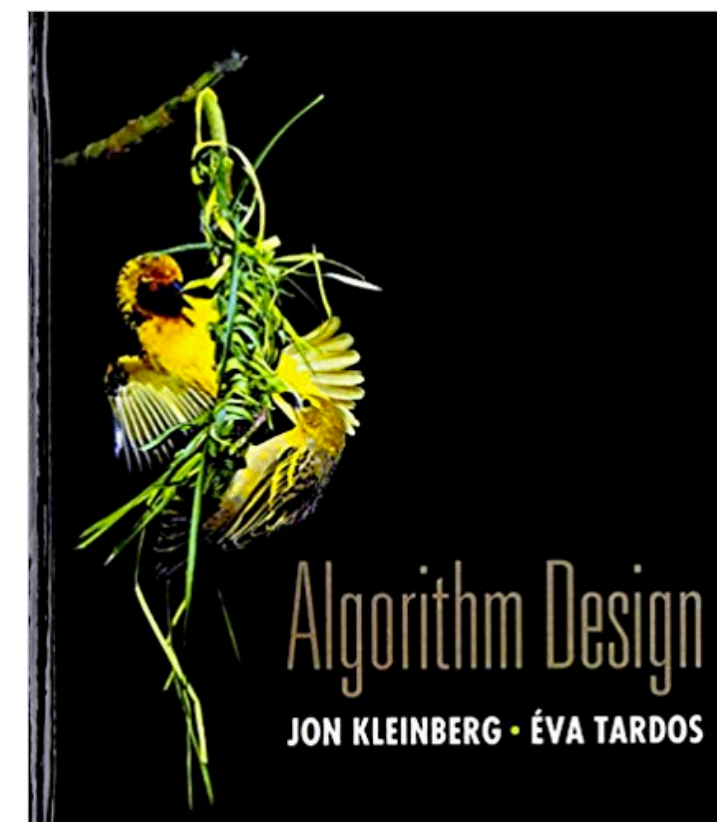
- **Instructor:** Prof. Fang Song.
- **Email:** fsong@pdx.edu.
- **Office hours:** F 8:30 - 10am and by appointment via Zoom.
- **TA:** Steven Libby (slibby@pdx.edu).
- **Zoom links:** “PSU Classes” Calendar.

- **Recommended texts**

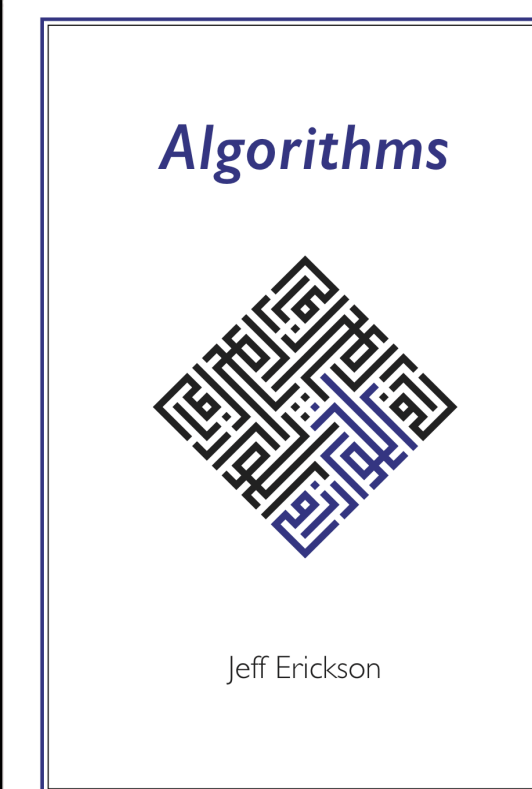
- Primary CLRS: E-copy available via PSU library



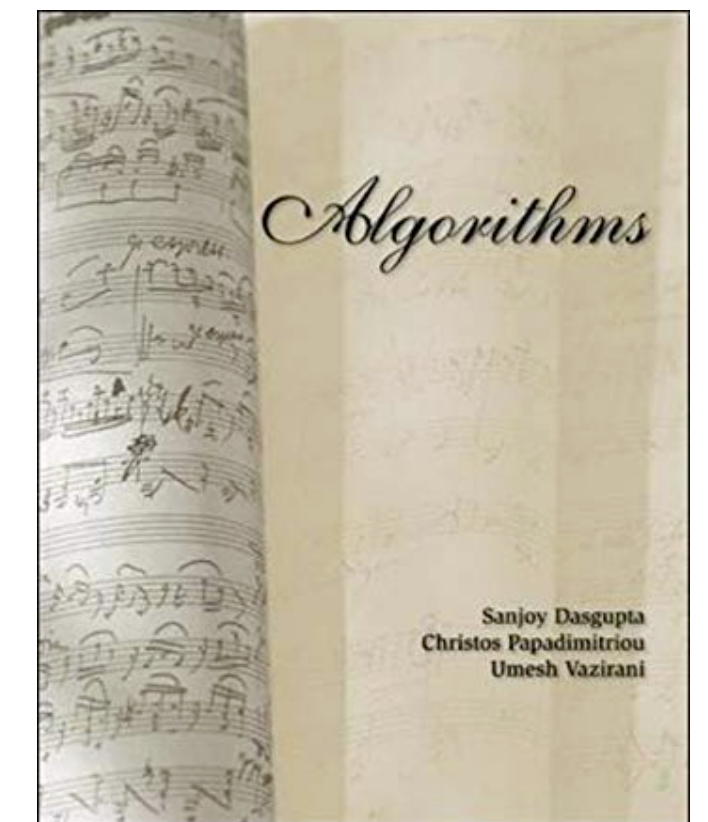
CLRS



KT



JE



DPV

Prerequisite & main topics

© CS 350 or equivalent

- Basic data structures and algorithms: sorting, graph traverse
- Math maturity: basics of combinatorics, linear algebra, probability

Comfortable with **READING** & **WRITING** mathematical **Proofs**

- Uncertain? Talk with me. Not a good idea to proceed if not ready.

© This course: more formal treatment and advanced materials

- Standard + selected topics at the end (quantum algorithms, etc.)

Policy

- ◎ **Final exam: 30%.**
- ◎ **Midterm exam: 25% week 5.**
- ◎ **Participation: 5%.**

- ◎ **Homework: 40%.**
 - Weekly.
 - Collaboration.
 - Find more on syllabus.

Policy cont'd

© Academic Integrity

- PSU Student Code of Conduct



© Academic accommodation

- Contact DRC (503-725-4150, drc@pdx.edu) and notify me.

© Recording lectures

- Comply with FERPA, the Acceptable Use Policy and PSU's Student Code of Conduct.
- Sharing outside this class not permitted.

How to succeed?

- ◎ Study the reading materials in advance.
- ◎ Start on assignment **EARLY!**
 - Make 20% progress per day for 5 days!
- ◎ Ask **a lot** of questions.
- ◎ Form **study groups**.

To-do #0: course webpage

You've probably accomplished it already. **Congrats!**

- ◎ Familiarize yourself with it https://fangsong.info/teaching/w21_5684_alg/
- ◎ “Schedule” page
 - Zoom links, reading materials, assignments.
- ◎ “Resource” page contains additional materials.

Check **regularly!**

To-do #1: get to know each other

- © It's very helpful to form a study group.
- © Gather town: <https://gather.town/app/CsidNQ8umdbRglZG/psu-w21-cs5684>

To-do #2: make a choice

- 1. Homework management**
 - a. Google classroom**
 - b. Gradescope**

- 2. Communication & discussion**
 - a. Email (via Google Groups)**
 - b. Slack**
 - c. Piazza or alternatives**

Complete the survey by Friday at <https://forms.gle/e56YsGkRnaiBHbYA6>

Now the real meat

- 1. Overview of algorithm design techniques.**
- 2. Overview of algorithm analysis.**
- 3. Growth of functions, asymptotic notations.**
- 4. LaTeX tutorial.**

Algorithmic techniques

★ **Reduction**: a meta technique, always keep in mind.

- ◎ **Brute force**
- ◎ **Divide and conquer**
 - Decompose a problem into smaller sub-problems and compose the solutions.
- ◎ **Dynamic programming**
 - Memorize soln's to subproblems that occur repeatedly.
- ◎ **Greediness**
 - Make a local optimal choice for subproblems.
- ◎ **Randomization**
- ◎ **... Creativity**

Algorithm analysis

◎ Correctness

- Pre-condition and post-conditions for each procedure (esp. **recursive** calls).
- Loop **invariant** for **iterative** algorithms.
- **Termination** in finite steps. E.x. decreasing of non-negative measure.

◎ Resource analysis

- Recurrences. $T(n) = 2T(n/2) + 3n$. Recursion tree, Master theorem.
- Amortized analysis, probabilistic analysis, ...
- Experimentation.

◎ Model of computation

- E.x. random-access machine (RAM). Unit cost per instruction and memory access.

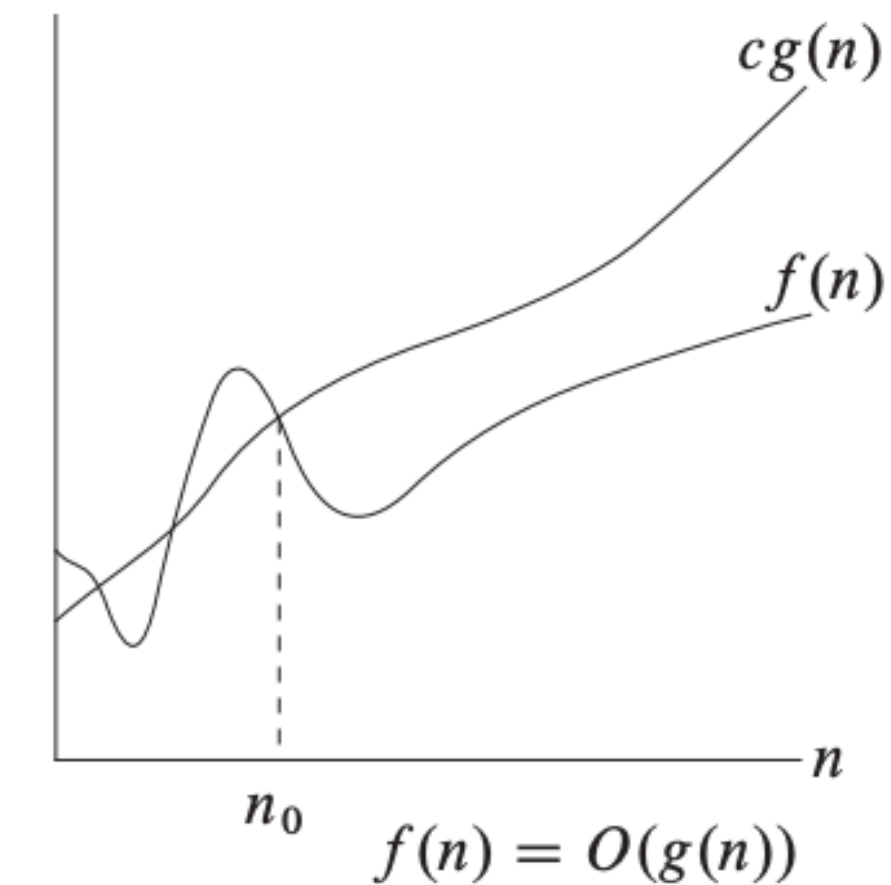
Asymptotic notations

◎ $O(\cdot)$, $\Omega(\cdot)$, $\Theta(\cdot)$, $o(\cdot)$, $\omega(\cdot)$

- Measure algorithm behaviors (by functions on integers) as problem size grows.
- Usually a good indicator of which alg. is preferable (except for small inputs).

◎ Defining $O(\cdot)$: asymptotic upper bound

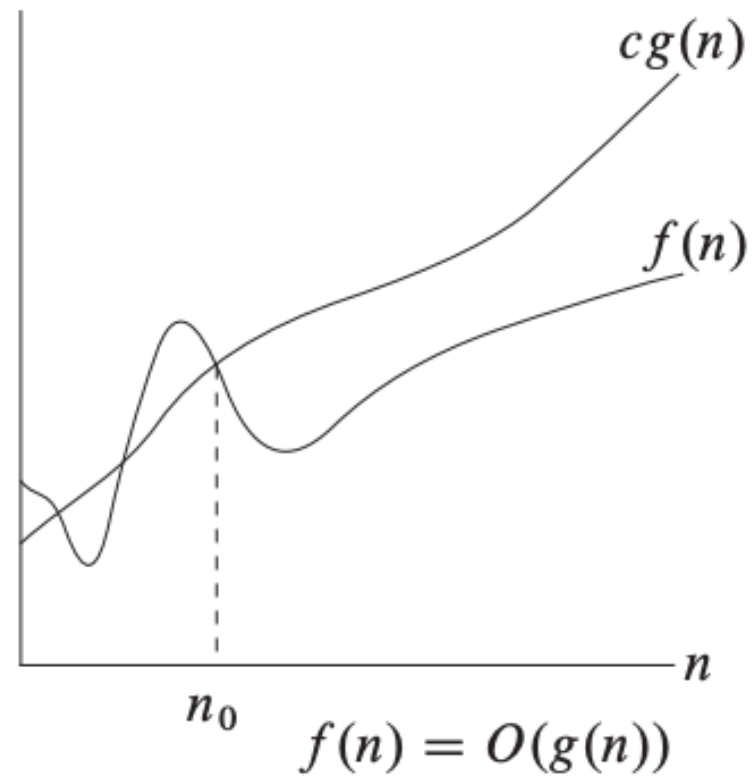
We write $f(n) = O(g(n))$ if there exist constants $c > 0, n_0 > 0$, such that $0 \leq f(n) \leq cg(n)$ for all $n \geq n_0$.



◎ $O(g(n))$ as a set

$$O(g(n)) := \{f(n) : \exists c > 0, n_0 > 0, \text{ such that } 0 \leq f(n) \leq c \cdot g(n) \text{ for all } n \geq n_0\}$$

Examples



$f(n) = O(g(n))$ if there **exist** constants $c > 0, n_0 > 0$, such that $0 \leq f(n) \leq cg(n)$ **for all** $n \geq n_0$.

⊙ $2n^2 = O(n^3)$

- $c = 1, n_0 = 2$.

- I.e., $2n^2 \in O(n^3)$

⊙ $f(n) = n^3 + O(n^2)$

- Meaning $f(n) = n^3 + h(n)$ for some $h(n) \in O(n^2)$

Exercise: sort by asymptotic order of growth

- | | |
|---------------|--------------------|
| 1. $n \log n$ | 6. n |
| 2. \sqrt{n} | 7. $n!$ |
| 3. $\log n$ | 8. $n^{1,000,000}$ |
| 4. n^2 | 9. $n^{1/\log n}$ |
| 5. 2^n | 10. $\log(n!)$ |

List them in **ascending** order: if f appears before g , then $f = O(g)$

9, 3, 2, 6, 1=10, 4, 8, 5, 7



Summary

Notation	... means ...	Think...	E.g.	Lim $f(n)/g(n)$
$f(n)=O(n)$	$\exists c>0, n_0>0, \forall n > n_0 : 0 \leq f(n) < cg(n)$	Upper bound	$100n^2 = O(n^3)$	If it exists, it is $< \infty$
$f(n)=\Omega(g(n))$	$\exists c>0, n_0>0, \forall n > n_0 : 0 \leq cg(n) < f(n)$	Lower bound	$n^{100} = \Omega(2^n)$	If it exists, it is > 0
$f(n)=\Theta(g(n))$	both of the above: $f=\Omega(g)$ and $f=O(g)$	Tight bound	$\log(n!) = \Theta(n \log n)$	If it exists, it is > 0 and $< \infty$
$f(n)=o(g(n))$	$\forall c>0, n_0>0, \forall n > n_0 : 0 \leq f(n) < cg(n)$	Strict upper bound	$n^2 = o(2^n)$	Limit exists, $=0$
$f(n)=\omega(g(n))$	$\forall c>0, n_0>0, \forall n > n_0 : 0 \leq cg(n) < f(n)$	Strict lower bound	$n^2 = \omega(\log n)$	Limit exists, $=\infty$

LaTeX

- ◎ Start with Overleaf.com
- ◎ Questions? Turn to Google and tex.stackexchange.com
- ◎ Short math guide for LaTeX
- ◎ Getting more serious?
 - A good text editor and packages/plugins, e.g., Emacs + AUCTeX
 - Version-control: GitHub
- ★ These things take time to learn, but you will not regret the effort!
- ★ Scrutinize your PDFs to see if anything could be nicer-looking.
- ★ Take pity on your poor reader.

LaTeX Live Demo

- ◎ What, why?
- ◎ Homework template file
 - Preamble,
 - Basic formatting: italics, bold, lists, reference, comments, table, etc.
 - Soln environment
- ◎ Math
 - Good practices

