



**S'20 CS410/510**

**Intro to  
quantum computing**

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## **Week 8**

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- Mixed states, density matrices
- General quantum operations
- POVM

# Exercise

1. Let  $I$  be identity on  $n$  qubits. Show that  $I = \sum_{x \in \{0,1\}^n} |x\rangle\langle x|$ .

$$\begin{aligned} |0\rangle\langle 0| &= \begin{pmatrix} 1 & 0 \\ 0 & 0 \end{pmatrix} \\ |1\rangle\langle 1| &= \begin{pmatrix} 0 & 0 \\ 0 & 1 \end{pmatrix} \end{aligned} \quad + = \begin{pmatrix} 1 & 0 \\ 0 & 1 \end{pmatrix}$$

$$|x\rangle\langle x| = \begin{pmatrix} \text{---} & & & \\ & \text{---} & & \\ & & \text{---} & \\ & & & \text{---} \end{pmatrix} \sim \mathbb{1}^{x \text{th}}$$

2. Let  $|A\rangle, |B\rangle$  be as defined below. Show that  $I = a|A\rangle\langle A| + b|B\rangle\langle B|$

$$a|A\rangle\langle A| = \left( \frac{1}{\sqrt{a}} \sum_{x \in A} |x\rangle \right) \left( \frac{1}{\sqrt{a}} \sum_{x'} \langle x'| \right)$$

- $A \subseteq \{0,1\}^n, B = \{0,1\}^n \setminus A$

- $|A\rangle := \frac{1}{\sqrt{a}} \sum_{x \in A} |x\rangle, |B\rangle := \frac{1}{\sqrt{b}} \sum_{x \in B} |x\rangle$

$$a = |A|$$

$$= \frac{1}{a} \sum_{\substack{x \in A \\ x' \in A}} |x\rangle\langle x'|$$

$$b|B\rangle\langle B| = \frac{1}{b} \sum_{\substack{x \in B \\ x' \in B}} |x\rangle\langle x'|$$

# Exercise

3. Let  $Z_f$  be as below. Show that  $Z_f = I - 2|A\rangle\langle A|$ . What is  $Z_f|A\rangle$ ?  $|B\rangle$

$$\bullet Z_f : |x\rangle \mapsto \begin{cases} -|x\rangle, & x \in A \\ |x\rangle, & x \notin A \end{cases}$$

$$Z_f |A\rangle = (I - 2|A\rangle\langle A|) |A\rangle$$

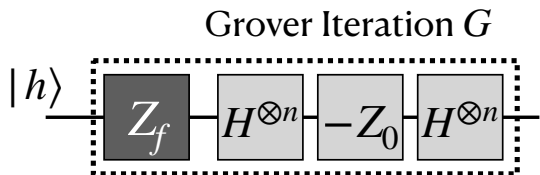
$$\stackrel{|B\rangle}{=} |A\rangle - 2 \underbrace{|A\rangle\langle A|A\rangle}_{\substack{1 \\ 1}} \underbrace{\langle A|A\rangle}_{\substack{1 \\ 0}} \stackrel{|A|B\rangle}{=} -|A\rangle$$

4. Let  $Z_0$  be as below. Show that  $Z_0 = I - 2|0^n\rangle\langle 0^n|$ .

$$\bullet Z_0 : |x\rangle \mapsto \begin{cases} -|x\rangle, & x = 0^n \\ |x\rangle, & x \neq 0^n \end{cases}$$

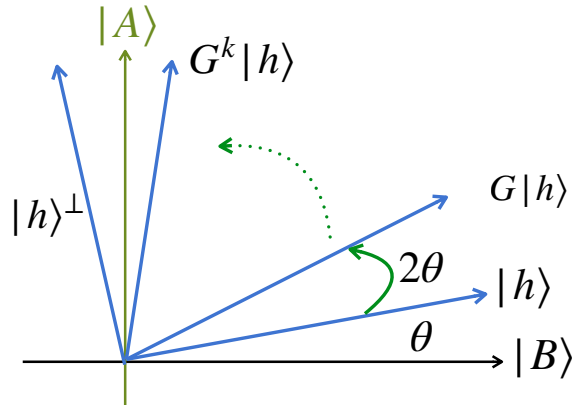
5. What is  $H^{\otimes n} Z_0 H^{\otimes n}$ ?

# Review: Grover's algorithm



- $Z_f = I - 2|A\rangle\langle A|$ : reflection about  $|B\rangle$
- $-HZ_0H = 2|h\rangle\langle h| - I$ : reflection about  $|h\rangle$
- $G = (-HZ_0H)Z_f$ : rotation by  $2\theta$

- $|A\rangle := \frac{1}{\sqrt{a}} \sum_{x \in A} |x\rangle$ ,  $|B\rangle := \frac{1}{\sqrt{b}} \sum_{x \in B} |x\rangle$
- $|h\rangle := H^{\otimes n} |0^n\rangle$
- $|h\rangle^\perp$ : orthogonal to  $|h\rangle$  on  $\text{span}\{|A\rangle, |B\rangle\}$



# Quantum algorithms so far

	Problem	Deterministic	Randomized	Quantum	
Partial function	Deutsch	2	2	1	oracle model
	Deutsch-Josza	$2^n/2$	$O(n)$	1	
	Simon	$2^n/2$	$\sqrt{2^n}$	$O(n^2)$	
	Order-finding Factoring $N$ (Kitaev/Shor)	$2^{O((\log N)^{1/3}(\log \log N)^{2/3})}$		$(\log N)^3$	
Total function	Unstructured search (Grover)	$\Omega(2^n)$		$\Theta(\sqrt{2^n})$	oracle model

# **Quantum information theory**

# An coarse taxonomy

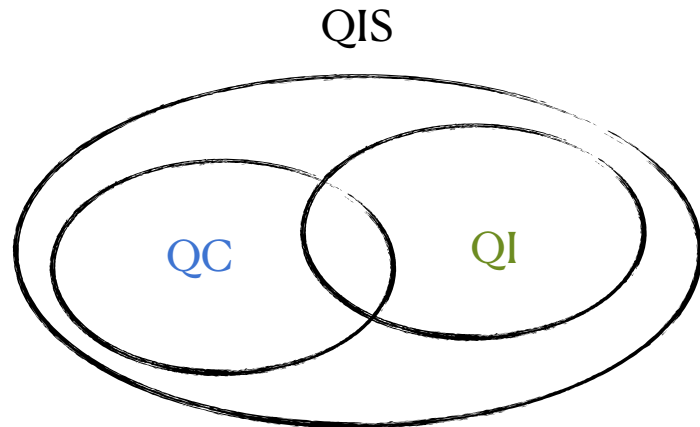
Quantum information science (QIS)

⊙ Quantum computing (QC): making information **useful**

- Algorithms, software, ...

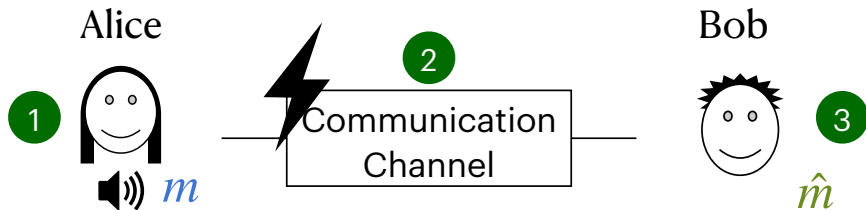
⊙ Quantum information (QI): making information **available**

- Elementary tasks: create, store, transmit, ...



# Basic communication scenario

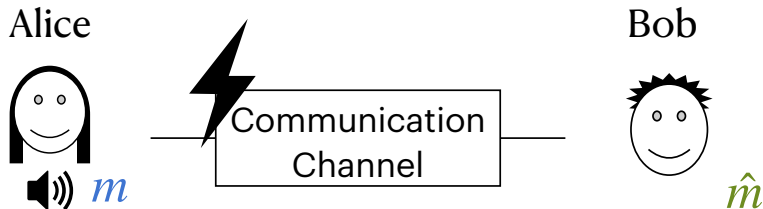
Goal: convey information from Alice to Bob



- 1 Alice: information source
- 2 Communication channel (resource): can you get everything I say in class?
- 3 Bob: because of noise, get disturbed  $\hat{m}$



# Central questions



0. What is information, mathematically?

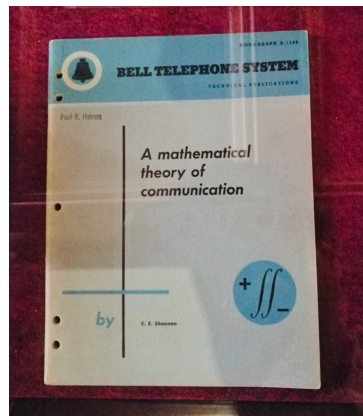
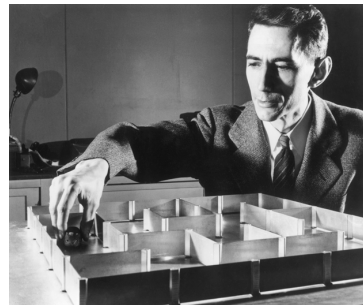
- Defining **bit** as unit of information

1. Assuming **noiseless** channel, how many bits needed to transmit  $m$ ?

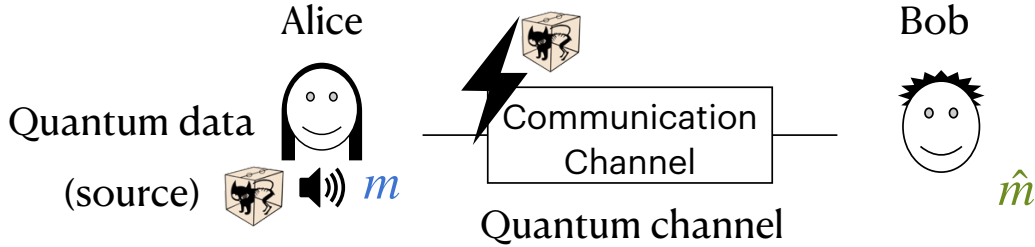
- Shannon **noiseless/source coding** theorem: **entropy**

2. Assuming **noisy** channel, how many bits can be transmitted **reliably**?

- Shannon **noisy-channel coding** theorem: **channel capacity**
- Tool: error correcting code



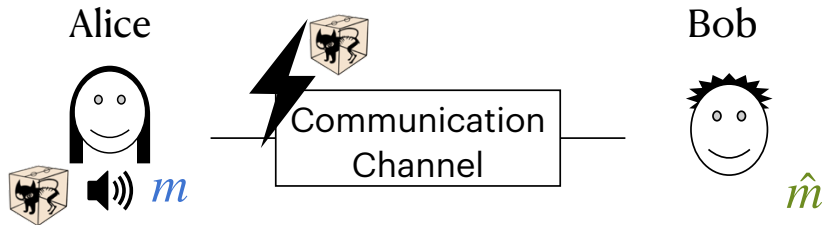
# The new quantum player



Channel \ Source	C	Q
Classical	<ol style="list-style-type: none"> <li>Shannon theory</li> </ol>	<ol style="list-style-type: none"> <li>Holevo's bound: # info. in qstates?</li> <li>Capacity to transmit C data</li> </ol>
Quantum	*teleportation	<ol style="list-style-type: none"> <li>Schumacher's Thm: compress Q data</li> <li>Quantum capacity</li> </ol>

- Noiseless channel
- Noisy channel

# The new quantum player



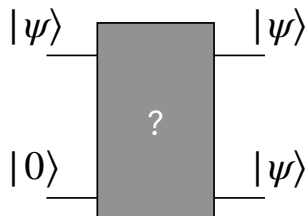
## © New resource: entanglement

- Teleportation, super dense coding
- Violation of Bell's inequality: validating quantum mechanics

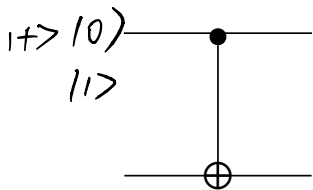
## © New challenges (easy for classical information)

- copying a quantum state?
- distinguishing states?

# Copy a quantum state?



● How about CNOT?

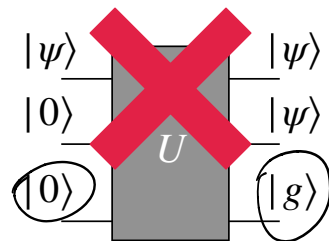


- $|0\rangle|0\rangle \mapsto |0\rangle|0\rangle$
- $|1\rangle|0\rangle \mapsto |1\rangle|1\rangle$
- $|+\rangle|0\rangle \mapsto |0\rangle|0\rangle + |1\rangle|1\rangle \neq |+\rangle|+\rangle$

$$(|0\rangle + |1\rangle)|0\rangle = |00\rangle + |10\rangle \xrightarrow{\text{CNOT}} |00\rangle + |11\rangle$$

# No-cloning theorem

**Theorem.** There is no valid quantum operation that maps an arbitrary (unknown) state  $|\psi\rangle$  to  $|\psi\rangle|\psi\rangle$ .



• **Proof. (Linearity)** Consider two states  $|\psi\rangle$  and  $|\psi'\rangle$

$$\begin{aligned} \bullet & \left( |\psi\rangle|0\rangle|0\rangle \right) \xrightarrow{U} |\psi\rangle|\psi\rangle|g\rangle \\ \bullet & \left( |\psi'\rangle|0\rangle|0\rangle \right) \xrightarrow{U} |\psi'\rangle|\psi'\rangle|g\rangle \end{aligned}$$

$U$  preserves inner product

$$\text{LHS } \langle 0| \langle 0| \langle \psi | \psi' \rangle |0\rangle |0\rangle \quad \langle g| \langle \psi | \langle \psi | \psi' \rangle |\psi\rangle |g\rangle \text{ RHS}$$

$$\langle \psi | \psi' \rangle \quad \langle \psi | \psi' \rangle \langle \psi | \psi' \rangle \langle g | g' \rangle$$

$$\langle \psi | \psi' \rangle (1 - \langle \psi | \psi' \rangle \langle g | g' \rangle) = 0$$

$$1 = \langle \psi | \psi' \rangle \langle g | g' \rangle$$

Two possibilities:  
 $\langle \psi | \psi' \rangle = 0$  or  $1$

# Density matrix formalism

# Another continent language

## State vector formalism

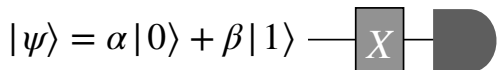
- State:  $|\psi\rangle \in \mathbb{C}^d$
- Unitary operation  $U : |\psi\rangle \mapsto U|\psi\rangle$
- Measuring in computational basis
  - $\sum_x \alpha_x |x\rangle$ : “x” w.p.  $|\alpha_x|^2$ , p.s.  $|x\rangle$

## Density matrix formalism

- State:  $\rho = |\psi\rangle\langle\psi|$  (density matrix)
  - Ex.  $|\psi\rangle = \alpha|0\rangle + \beta|1\rangle$   $|\psi\rangle\langle\psi| = \begin{pmatrix} |\alpha|^2 & \alpha^*\beta \\ \alpha\beta^* & |\beta|^2 \end{pmatrix}$
- Unitary  $U : \rho \mapsto U\rho U^\dagger$ 
  - $|\psi\rangle \xrightarrow{U} U|\psi\rangle$ ,  $(\underbrace{U|\psi\rangle}_{\rho})(\langle\psi|U^\dagger) = U\rho U^\dagger$
- Measuring in computational basis
  - $\rho = \sum_{x,x'} \alpha_x \alpha_{x'}^* |x\rangle\langle x'|$ : “x” w.p.  $\langle x|\rho|x\rangle$ , p.s.  $|x\rangle\langle x|$

# Exercise

1. Analyze the circuit below under both formalisms.



$\downarrow X$

$$\alpha|1\rangle + \beta|0\rangle$$

$\downarrow u$

$$0 \text{ w.p. } |\beta|^2 \quad |0\rangle$$

$$1 \text{ w.p. } |\alpha|^2 \quad |1\rangle$$

$$\rho = |\alpha|^2 |0\rangle\langle 0| + \alpha\beta^* |0\rangle\langle 1| + \alpha^*\beta |1\rangle\langle 0| + |\beta|^2 |1\rangle\langle 1|$$

$\downarrow X$

$$X\rho X^\dagger = |\alpha|^2 X(|0\rangle\langle 0|)X^\dagger + \dots$$

$$\rho' = |\alpha|^2 |1\rangle\langle 1| + \dots + \dots$$

$\downarrow u$  "1" w.p.  $\langle 1|\rho'|1\rangle = |\alpha|^2 \quad |1\rangle\langle 1|$

2. Consider two qubits in state  $|+\rangle|-\rangle$ . Write down its density matrix.

$$|+\rangle|-\rangle = \begin{pmatrix} 1 \\ 1 \end{pmatrix} \otimes \begin{pmatrix} 1 \\ -1 \end{pmatrix}$$

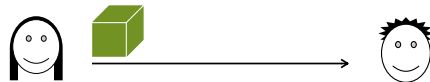
$$= \begin{pmatrix} 1 \\ -1 \\ 1 \\ -1 \end{pmatrix}$$

$$\rho = \begin{pmatrix} 1 \\ -1 \\ 1 \\ -1 \end{pmatrix} \begin{pmatrix} 1 & -1 & 1 & -1 \end{pmatrix}$$

$$= \begin{pmatrix} & & & \\ & & & \\ & & & \\ & & & \end{pmatrix}_{4 \times 4}$$



# Pure states vs. mixed states



- Alice flips a coin, prepare  $|0\rangle$  or  $|1\rangle$  accordingly.
- Bob receives the register (Alice's coin unknown). How to describe his state?
  - $\{(\frac{1}{2}, |0\rangle), (\frac{1}{2}, |1\rangle)\}$  no compact representation as state vectors
  - Density matrix representation:  $\rho = \frac{1}{2} |0\rangle\langle 0| + \frac{1}{2} |1\rangle\langle 1|$
- This is called a **mixed** state. In contrast,  $|\psi\rangle$  is called a **pure** state.

# Exercise

1. Alice flips a coin and prepares a qubit as follows. She then sends the qubit (but not the coin) to Bob. How to describe Bob's state?



HEADS:  $|+\rangle$

TAILS:  $|-\rangle$

2. Write down the density matrix explicitly and compare with the previous slide.

$$\rho = \frac{1}{2} |+\rangle\langle+| + \frac{1}{2} |-\rangle\langle-|$$

$$\rho = \frac{1}{2} |0\rangle\langle 0| + \frac{1}{2} |1\rangle\langle 1|$$

$$= \frac{1}{2} \begin{pmatrix} 1 & 0 \\ 0 & 1 \end{pmatrix}_{2 \times 2}$$

$$\rho^a = \frac{1}{4} \begin{pmatrix} 1 & 0 \\ 0 & 1 \end{pmatrix}$$

$$= \frac{1}{2} \begin{pmatrix} 1 & 0 \\ 0 & 1 \end{pmatrix}_{2 \times 2}$$

# General mixed states

● Mixed state = a probability distribution (mixture) over pure states

•  $\{(p_i, |\psi_i\rangle) : i = 1, \dots, k\} : \rho = \sum_i p_i |\psi_i\rangle\langle\psi_i|$

● Properties of density matrices

•  $\text{tr}(\rho) = 1$

•  $\rho$  is pure iff.  $\text{tr}(\rho^2) = 1$  (Think of examples in previous slides)

$\text{tr}[(|0\rangle\langle 0|)^2] = \text{tr}[|0\rangle\langle 0|] = 1$ ;       $\rho = |\psi\rangle\langle\psi|$ ;  $\rho^2 = \underbrace{(|\psi\rangle\langle\psi|)(|\psi\rangle\langle\psi|)}_{= \rho}$

•  $\rho$  is positive semi-definite, i.e.,  $\langle\psi|\rho|\psi\rangle \geq 0$ .

# Operations on mixed states

• Unitary  $U : \rho \mapsto U\rho U^\dagger$

$$\{ (p_i, |\psi_i\rangle) \} \quad \rho = \sum_i p_i |\psi_i\rangle\langle\psi_i|$$

$$U |\psi_i\rangle\langle\psi_i| U^\dagger$$

$\xrightarrow{U}$

$$= \sum_i p_i U |\psi_i\rangle\langle\psi_i| U^\dagger$$

$$= U \left( \sum_i p_i |\psi_i\rangle\langle\psi_i| \right) U^\dagger$$

$$= U \rho U^\dagger$$

• Measurement: "x" with prob.  $\langle x | \rho | x \rangle$

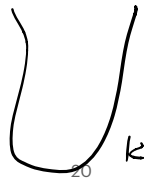
$\forall i:$

$$|\psi_i\rangle\langle\psi_i| : \text{"x"} \quad \langle x | |\psi_i\rangle\langle\psi_i| | x \rangle$$

i<sup>th</sup> bin:  $x \cdot p_x^i \rightarrow$  Question:



...



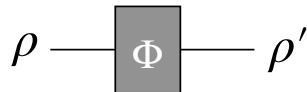
if pick i<sup>th</sup> bin w.p.  $p_i$   
& pick  $x$  from i<sup>th</sup> bin ( $p_x^i$ )

$$\text{then } P_r[x] = \sum_i p_i P_r[x | i]$$

$$\rightarrow P_r[\text{meas. } x] = \sum_i p_i \langle x | |\psi_i\rangle\langle\psi_i| | x \rangle$$

$$= \langle x | \left( \sum_i p_i |\psi_i\rangle\langle\psi_i| \right) | x \rangle$$

# General quantum operations



Let  $A_1, A_2, \dots, A_m$  be matrices satisfying  $\sum_{j=1}^m A_j^\dagger A_j = I$ .

Then the mapping  $\rho \mapsto \sum_{j=1}^m A_j \rho A_j^\dagger$  is a general quantum operator.

- N.B.  $A_i$  need NOT be square matrices
- Also known as **quantum channels**
  - admissible operations, completely positive trace preserving maps

# Examples of quantum channels

1. Unitary  $U^\dagger U = I: \rho \mapsto U\rho U^\dagger$

2. Decoherence channel  $A_0 = |0\rangle\langle 0|, A_1 = |1\rangle\langle 1|$

• Check validity:  $A_0^\dagger A_0 + A_1^\dagger A_1 = \mathbb{1}$

$$(|0\rangle\langle 0|)(|0\rangle\langle 0|) + (|1\rangle\langle 1|)(|1\rangle\langle 1|) = \mathbb{1}$$

• Apply to  $|\psi\rangle = \alpha|0\rangle + \beta|1\rangle$

$$\rho \xrightarrow{\Phi} \sum_i A_i \rho A_i^\dagger$$

$$A_0 (|\alpha|^2 |0\rangle\langle 0|) A_0^\dagger + A_1 (|\beta|^2 |1\rangle\langle 1|) A_1^\dagger$$

$$\rho = \begin{pmatrix} |\alpha|^2 & \alpha\beta^* \\ \alpha^*\beta & |\beta|^2 \end{pmatrix} = |\alpha|^2 |0\rangle\langle 0| + \alpha\beta^* |0\rangle\langle 1| + \alpha^*\beta |1\rangle\langle 0| + |\beta|^2 |1\rangle\langle 1|$$

• Compare to measurement:

$$A_0 \rho A_0^\dagger = |\alpha|^2 |0\rangle\langle 0|$$

$$A_1 \rho A_1^\dagger = |\beta|^2 |1\rangle\langle 1|$$

$$\Phi(\rho) = \begin{pmatrix} |\alpha|^2 & 0 \\ 0 & |\beta|^2 \end{pmatrix}$$

# Examples of quantum channels

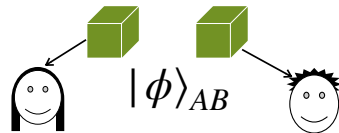
3. **Partial trace**  $A_0 = I \otimes \langle 0| = \begin{pmatrix} 1 & 0 & 0 & 0 \\ 0 & 0 & 1 & 0 \end{pmatrix}$ ,  $A_1 = I \otimes \langle 1| = \begin{pmatrix} 0 & 1 & 0 & 0 \\ 0 & 0 & 0 & 1 \end{pmatrix}$

- Check validity:

- Apply to  $|0\rangle\langle 0| \otimes |+\rangle\langle +|$

- Apply to  $|\phi\rangle = \frac{1}{\sqrt{2}}(|00\rangle + |11\rangle)$

# Exercise

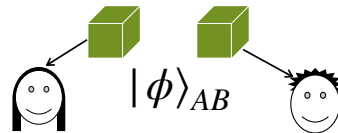


1. let  $Tr_B$  denote partial trace of subsystem  $B$ . Suppose Alice and Bob shares two qubits in state  $|\phi\rangle_{AB}$ .

- Apply  $Tr_B$  to  $|\phi\rangle_{AB} = \frac{1}{\sqrt{2}}(|00\rangle + |11\rangle)$
- Apply  $Tr_B$  to  $|\phi\rangle_{AB} = \frac{1}{\sqrt{2}}(|01\rangle + |10\rangle)$
- Is Alice able to tell the two cases on her side?



# Exercise



2. let  $Tr_B$  denote partial trace of subsystem  $B$ . Suppose Alice and Bob shares two qubits in state  $|\phi\rangle_{AB}$ .

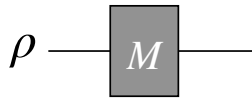
- Apply  $Tr_B$  to  $|\phi\rangle_{AB} = \frac{3}{5}|00\rangle + \frac{4}{5}|11\rangle$

- Apply  $Tr_B$  to  $|\phi\rangle_{AB} = \frac{4}{5}|00\rangle - \frac{3}{5}|11\rangle$

- Is Alice able to tell the two cases on her side?

# General measurement

- A **measurement** is described by a collection of matrices  $M = \{M_a : a \in \Gamma\}$  with possible outcomes  $\Gamma$  satisfying  $\sum_{a \in \Gamma} M_a^\dagger M_a = I$ .



outcome	probability	posterior state
$a$	$\text{Tr}(M_a \rho M_a^\dagger)$	$\frac{M_a \rho M_a^\dagger}{\text{Tr}(M_a \rho M_a^\dagger)}$

- Example.  $M_0 = |0\rangle\langle 0| \otimes I, M_1 = |1\rangle\langle 1| \otimes I, \Gamma = \{0, 1\}$ .

- Measure  $|\psi\rangle = |+\rangle|0\rangle$

$$|\psi\rangle\langle\psi| = |+\rangle\langle+| \otimes |0\rangle\langle 0|$$

$$M_0 (|+\rangle\langle+| \otimes |0\rangle\langle 0|) M_0^\dagger$$

$$= (|0\rangle\langle 0| + |1\rangle\langle 1|) \otimes |0\rangle\langle 0| = \frac{1}{2} |0\rangle\langle 0| \otimes |0\rangle\langle 0| = \begin{pmatrix} \frac{1}{2} & 0 & 0 & 0 \\ 0 & 0 & 0 & 0 \\ 0 & 0 & 0 & 0 \\ 0 & 0 & 0 & 0 \end{pmatrix}$$

# Projective measurement & POVM

- Projective (von Neumann) measurement:  $M_a$  projections ( $M_a^2 = M_a$ ).
  - Complete projective measurement  $M_a = |\psi_a\rangle\langle\psi_a|$  and  $\{|\psi_a\rangle\}$  an orthonormal basis
  - $\equiv$  measurement under basis  $\{|\psi_a\rangle\}$
- Positive-operator-valued measurement (POVM) measurement
  - $\Pr[a] = \text{Tr}(M_a \rho M_a^\dagger) = \text{Tr}(\overset{E_a}{M_a^\dagger M_a} \rho)$        $\text{Tr}(AB) = \text{Tr}(BA)$
  - Suffice to specify POVM elements  $\{E_a = M_a^\dagger M_a : a \in \Gamma\}$

# Logistics

● HW6 due next Sunday

● Project

- Week10 office hour: slots available
- Presentations
  - Pre-record your talk by zoom, powerpoint, ... Keep it 20 - 25 mins
  - Live Q&A in class
  - Participate in all talks and fill out peer-evaluation

# Discussion on Google's experiment

