

F, 04/10/2020



Portland State U

S'20 CS 410/510

**Intro to
quantum computing**

Fang Song

Week 2

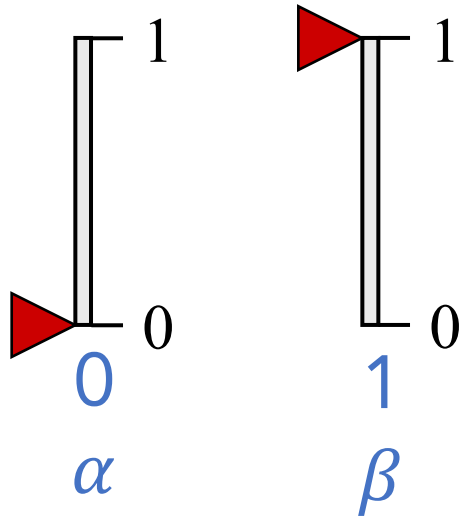
- Multiple qubits, tensor product
- Quantum circuit model
- Quantum superdense coding
- Quantum teleportation

Credit: based on slides by Richard Cleve

Logistics

- HW1 due Sunday
 - Work in groups, write up individually
- Project
 - Form groups of 2-3 persons by next week
- Workflow
 - Work on **pre-class** materials: 80% success depends on it!
 - In-class: practice what you studied and extend to new topics
 - Post-class: review and reinforce

Review: qubit



Superposition

- Amplitudes $\alpha, \beta \in \mathbb{C}, |\alpha|^2 + |\beta|^2 = 1$
- Explicit state is $\begin{pmatrix} \alpha \\ \beta \end{pmatrix} \in \mathbb{C}^2$
(2-norm / Euclidean norm = 1)
- **Cannot** explicitly extract α and β
(only statistical inference)

Dirac bra/ket notation

- **Ket:** $|\psi\rangle$ always denotes a column vector

Convention: $|0\rangle = \begin{pmatrix} 1 \\ 0 \end{pmatrix}$, $|1\rangle = \begin{pmatrix} 0 \\ 1 \end{pmatrix}$

Ex. $|\psi\rangle = \begin{pmatrix} \alpha_1 \\ \alpha_2 \\ \vdots \\ \alpha_d \end{pmatrix}$

- **Bra:** $\langle\psi|$ always denotes a row vector that is the conjugate transpose of $|\psi\rangle$

Ex. $\langle\psi| = (\alpha_1^*, \alpha_2^*, \dots, \alpha_d^*)$

- **Inner product:** $\langle\psi|\phi\rangle$ denotes $\langle\psi| \cdot |\phi\rangle$

- Vectors to **scalar**

Ex. $\langle 0|1\rangle = (1 \ 0) \begin{pmatrix} 0 \\ 1 \end{pmatrix} = 0$

- **Outer product:** $|\psi\rangle\langle\phi|$ denotes $|\psi\rangle \cdot \langle\phi|$

- Vectors to **matrix**

Ex. $|0\rangle\langle 1| = \begin{pmatrix} 1 \\ 0 \end{pmatrix} (0 \ 1) = \begin{pmatrix} 0 & 1 \\ 0 & 0 \end{pmatrix}$

Basic operations on a qubit

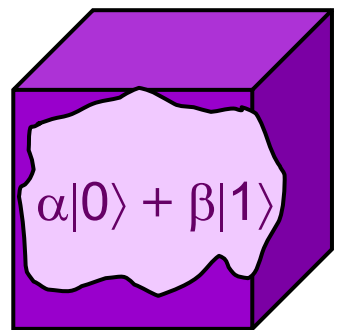
0. Initialize qubit to $|0\rangle$ or to $|1\rangle$
1. Apply a **unitary** operation U ($U^\dagger U = I$)

$$X = \begin{pmatrix} 0 & 1 \\ 1 & 0 \end{pmatrix}, \quad X|0\rangle = |1\rangle, X|1\rangle = |0\rangle$$
$$Z = \begin{pmatrix} 1 & 0 \\ 0 & -1 \end{pmatrix}, \quad Z|0\rangle = |0\rangle, Z|1\rangle = -|1\rangle$$
$$H = \frac{1}{\sqrt{2}} \begin{pmatrix} 1 & 1 \\ 1 & -1 \end{pmatrix}, \quad H|0\rangle = |+\rangle, H|1\rangle = |-\rangle$$

Linear map $A \leftrightarrow$ **matrix** A
Apply A to $|\psi\rangle \leftrightarrow$ matrix mult. $A|\psi\rangle$

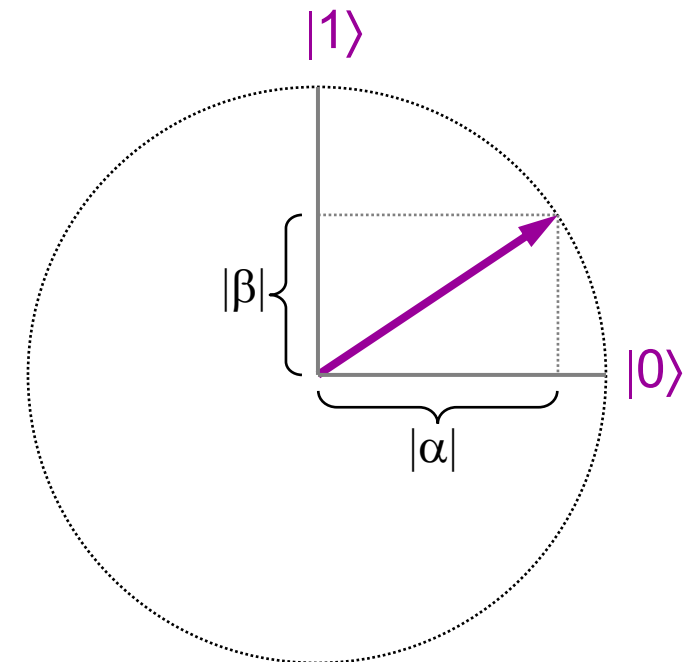
Basic operations on a qubit

0. Initialize qubit to $|0\rangle$ or to $|1\rangle$
1. Apply a **unitary** operation U ($U^\dagger U = I$)
2. Perform a “standard” measurement:



$$\mapsto \begin{cases} 0 & \text{with prob } |\alpha|^2 & |0\rangle \\ 1 & \text{with prob } |\beta|^2 & |1\rangle \end{cases} \quad \text{posterior state}$$

... and the quantum state collapses



N.B. There exist other quantum operations, but they can all be “**simulated**” by the aforementioned types

A few tips


- **Linearity.** Let A be a linear map. Any $v_i \in \mathbb{C}^d, c_i \in \mathbb{C}, i = 1, \dots, k$

$$A\left(\sum_i c_i \cdot v_i\right) = \sum_i c_i \cdot A(v_i)$$

→ A is **uniquely determined** by its action on a basis

- Let $u_1, \dots, u_d \in \mathbb{C}^d$ be a basis → $\forall v \in \mathbb{C}^d$ can be expressed by $v = \sum_i c_i u_i$
 - Given $A(u_i) = w_i, i = 1, \dots, d$ → $Av = A(\sum_i c_i u_i) = \sum_i c_i A(u_i) = \sum_i c_i w_i$
- When Dirac notation unclear, convert to vectors/matrices
 - When Dirac notation unclear, convert to vectors/matrices

Two qubits: composed system


$$\underbrace{\alpha|0\rangle + \beta|1\rangle} \otimes \underbrace{\alpha'|0\rangle + \beta'|1\rangle} = \alpha\alpha'|00\rangle + \alpha\beta'|01\rangle + \beta\alpha'|10\rangle + \beta\beta'|11\rangle$$

▪ Tensor product \otimes

$$[A]_{m \times n} \otimes [B]_{k \times \ell} = \begin{pmatrix} a_{11}B & a_{12}B & \dots & a_{1n}B \\ a_{21}B & a_{22}B & \dots & a_{2n}B \\ \vdots & \vdots & \dots & \dots \\ a_{m1}B & a_{m2}B & \dots & a_{mn}B \end{pmatrix}_{mk \times n\ell}$$

Two qubits: composed system

$$[A]_{m \times n} \otimes [B]_{k \times \ell} = \begin{pmatrix} a_{11}B & a_{12}B & \dots & a_{1n}B \\ a_{21}B & a_{22}B & \dots & a_{2n}B \\ \vdots & \vdots & \dots & \dots \\ a_{m1}B & a_{m2}B & \dots & a_{mn}B \end{pmatrix}_{mk \times n\ell}$$

$$\text{Ex. } |00\rangle := |0\rangle \otimes |0\rangle = \begin{pmatrix} 1 \\ 0 \end{pmatrix} \otimes \begin{pmatrix} 1 \\ 0 \end{pmatrix} = \begin{pmatrix} 1 \cdot \begin{pmatrix} 1 \\ 0 \end{pmatrix} \\ 0 \cdot \begin{pmatrix} 1 \\ 0 \end{pmatrix} \end{pmatrix} = \begin{pmatrix} 1 \\ 0 \\ 0 \\ 0 \end{pmatrix}$$

$$|01\rangle = \begin{pmatrix} 0 \\ 0 \\ 0 \\ 0 \end{pmatrix}, |10\rangle = \begin{pmatrix} 0 \\ 0 \\ 1 \\ 0 \end{pmatrix}, |11\rangle = \begin{pmatrix} 0 \\ 0 \\ 0 \\ 1 \end{pmatrix} \quad |\psi\rangle|\phi\rangle := |\psi\rangle \otimes |\phi\rangle$$

General n -qubit systems

- Probabilistic states

$$\forall x \in \{0,1\}^n, p_x \geq 0$$
$$\sum_x p_x = 1$$
$$\begin{pmatrix} p_{000} \\ p_{001} \\ p_{010} \\ p_{011} \\ p_{100} \\ p_{101} \\ p_{110} \\ p_{111} \end{pmatrix}$$

- Quantum states

$$\forall x \in \{0,1\}^n, \alpha_x \in \mathbb{C}$$
$$\sum_x |\alpha_x|^2 = 1$$
$$\begin{pmatrix} \alpha_{000} \\ \alpha_{001} \\ \alpha_{010} \\ \alpha_{011} \\ \alpha_{100} \\ \alpha_{101} \\ \alpha_{110} \\ \alpha_{111} \end{pmatrix}$$

Dirac notation: $|000\rangle, |001\rangle, |010\rangle, \dots, |111\rangle$ are **basis** vectors

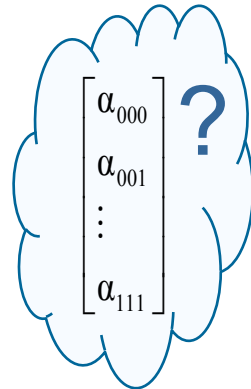
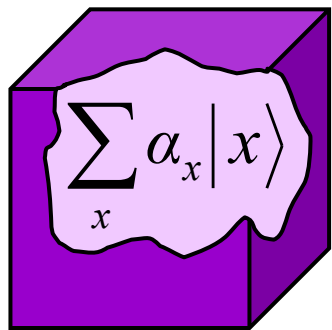
→ Any state can be written as $|\psi\rangle = \sum_x \alpha_x |x\rangle$

Operations on n -qubit states

- Unitary operations:
($U^\dagger U = I$)

$$\sum_x \alpha_x |x\rangle \mapsto U \left(\sum_x \alpha_x |x\rangle \right)$$

- Measurements:



		posterior state	
{	000	with prob $ \alpha_{000} ^2$	$ 000\rangle$
	001	with prob $ \alpha_{001} ^2$	$ 001\rangle$
	\vdots	\vdots	\vdots
	111	with prob $ \alpha_{111} ^2$	$ 111\rangle$

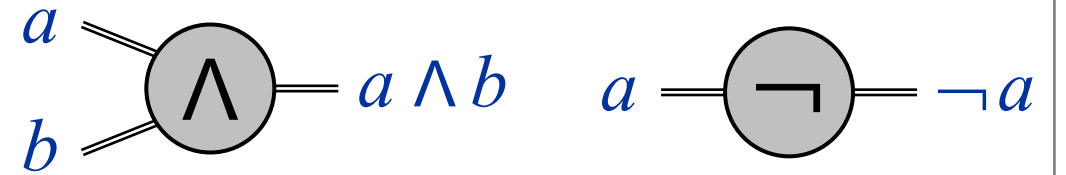
... and the quantum state collapses

Model of computation

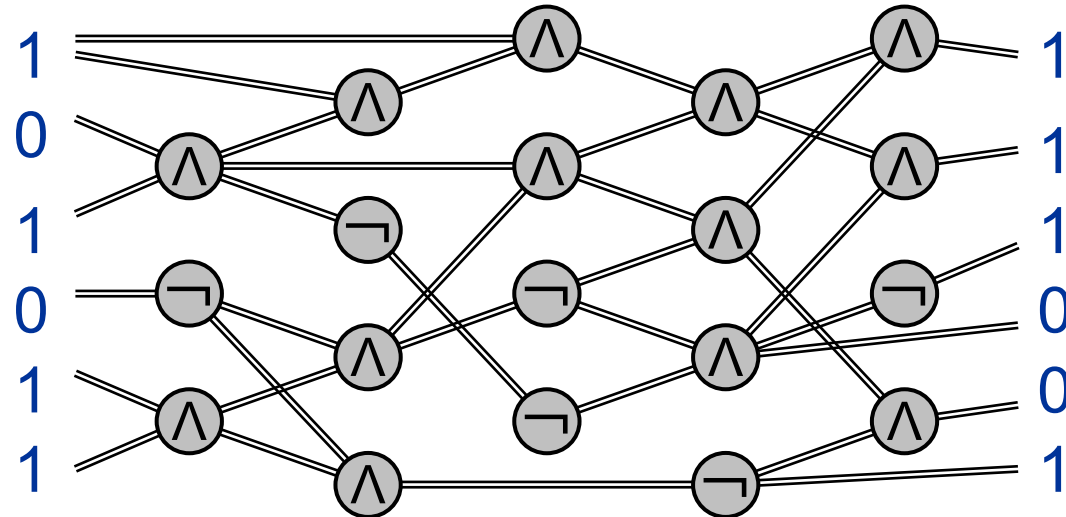
Classical Boolean circuits

Bit 0/1

Data flow \longrightarrow



Classical circuits:



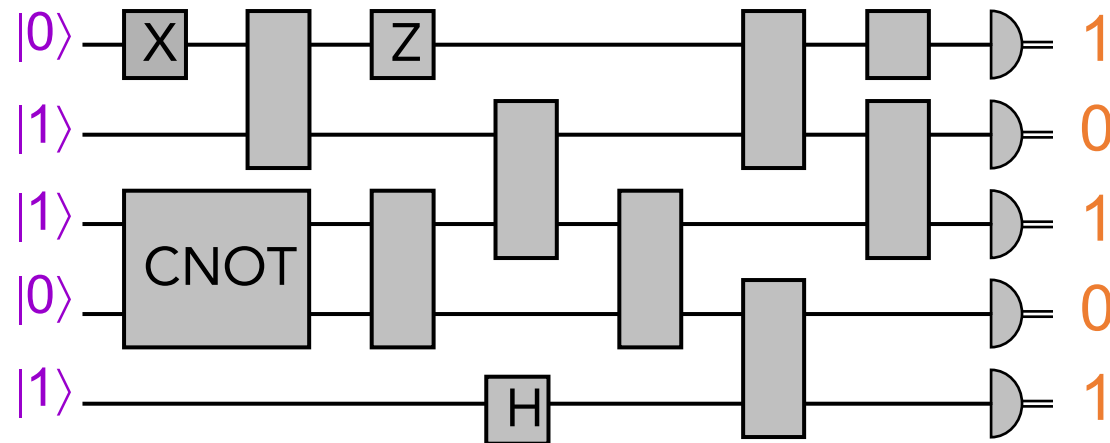
Quantum circuit model

Qubit

$$|\psi\rangle = \alpha|0\rangle + \beta|1\rangle$$

Data flow 

Quantum circuits:

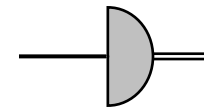


$$|b\rangle \xrightarrow{X} |\neg b\rangle$$

$$|a\rangle \xrightarrow{\text{Control}} |a\rangle$$

$$|b\rangle \xrightarrow{\text{Target}} |a \oplus b\rangle$$

$$|a\rangle|b\rangle \mapsto |a\rangle|a \oplus b\rangle$$



(standard) Measure

The power of computation

- **Computability:** can you solve it, in principle?

[Given program code, will this program terminate or loop indefinitely?]

Uncomputable!

Church-Turing Thesis. A problem can be computed in any *reasonable* model of computation *iff.* it is computable by a **Boolean circuit**.

- **Complexity:** can you solve it, under resource constraints?

[Can you factor a 1024-bit integer in 3 seconds?]

Extended Church-Turing Thesis. A function can be computed *efficiently* in any *reasonable* model of computation *iff.* it is efficiently computable by a **Boolean circuit**.



Product state vs. entangled state

- Product state $|\psi\rangle_{AB} = |\psi\rangle_A \otimes |\phi\rangle_B$

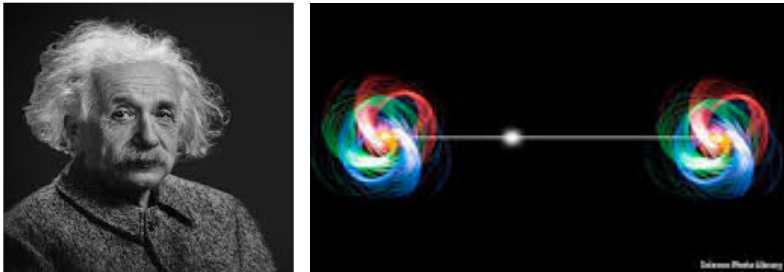
$$|\psi\rangle_{AB} = \overbrace{\alpha|0\rangle + \beta|1\rangle}^{\text{A}} \otimes \overbrace{\alpha'|0\rangle + \beta'|1\rangle}^{\text{B}}$$

- $|\psi\rangle_{AB}$ an arbitrary 2-qubit state:
Can we always write it as $|\psi\rangle_A \otimes |\phi\rangle_B$ for some $|\psi\rangle_A$ and $|\phi\rangle_B$?

Product state vs. entangled state

- **Entangled** state: $|\psi\rangle_{AB} \neq |\psi\rangle_A \otimes |\psi\rangle_B$ for any $|\psi\rangle_A$ and $|\phi\rangle_B$

Ex. $|\Phi\rangle = \frac{1}{\sqrt{2}} (|00\rangle + |11\rangle)$ EPR (Einstein–Podolsky–Rosen) pair



- Mathematically, not surprising:
A & B **correlated**
- Physically, non-classical correlation,
“spooky” action at a distance

- Cor. need to speak of **state of entire system** than individuals

Exercise: correlation & entanglement

1. Consider two bits a & b whose joint state (i.e., prob. distribution) is

described by probabilistic vector $v = \begin{pmatrix} 1/2 \\ 0 \\ 0 \\ 1/2 \end{pmatrix}$.

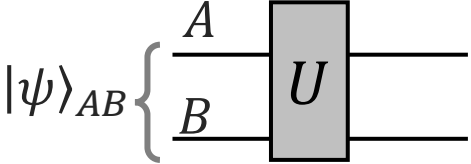
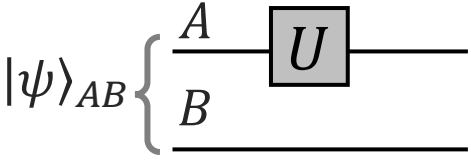
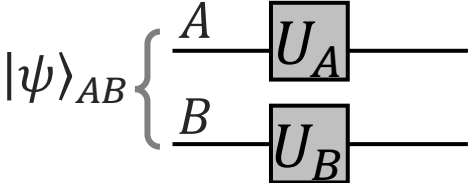
- What is the probability that $ab = 11$?
- Does there exist two-dimensional probabilistic vectors u_A and u_B such that $v = u_A \otimes u_B$?

Exercise: correlation & entanglement

2. Prove that the EPR state $|\Phi\rangle = \frac{1}{\sqrt{2}}(|00\rangle + |11\rangle)$ cannot be written as $|\psi\rangle \otimes |\phi\rangle$ for any choice of $|\psi\rangle, |\phi\rangle \in \mathbb{C}^2$.

Two-qubit gates

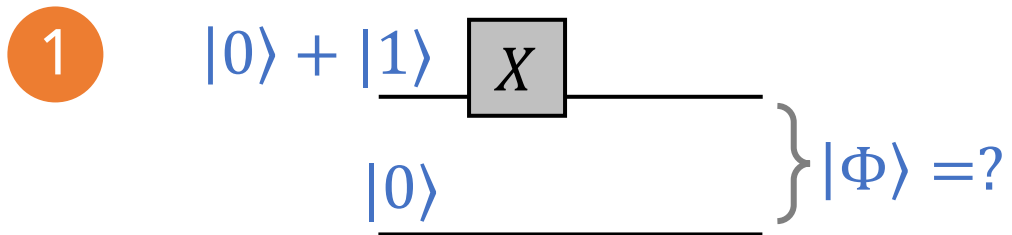
Given two qubits in state $|\psi\rangle_{AB}$

Description	Algebra	Circuit
Apply unitary U to $ \psi\rangle_{AB}$	$U \psi\rangle_{AB}$	
Apply unitary U to qubit A	$U \otimes I \psi\rangle_{AB}$	
Apply unitary U_A to qubit A & unitary U_B to qubit B	$U_A \otimes U_B \psi\rangle_{AB}$	

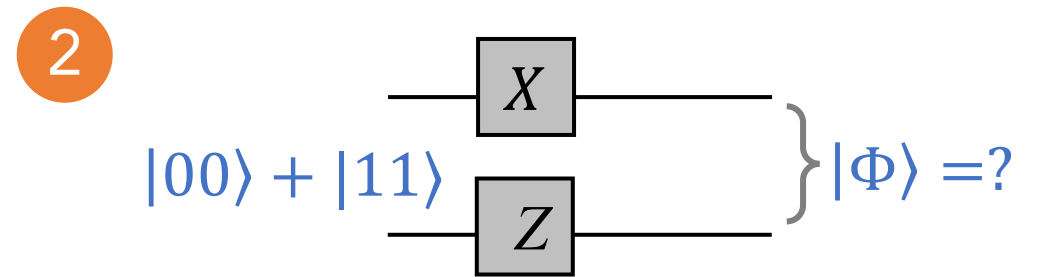
■ Facts

- Given unitary U, V , $U \otimes V$ is also unitary.
- $(U \otimes V)(A \otimes B) = UA \otimes VB$

Exercise: two-qubit gates



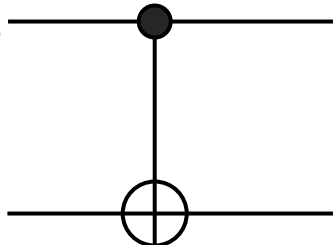
i.e. $|\Phi\rangle = X \otimes I((|0\rangle + |1\rangle)_A \otimes |0\rangle_B)$
 $= ?$



i.e. $|\Phi\rangle = X \otimes Z(|00\rangle + |11\rangle)$
 $= ?$

Exercise: CNOT

Control $|a\rangle$



Target $|b\rangle$

$|a\rangle$

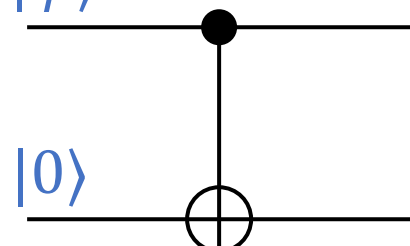
$|a \oplus b\rangle$

CNOT: $|00\rangle \mapsto |00\rangle$
 $|01\rangle \mapsto |00\rangle$
 $|10\rangle \mapsto |11\rangle$
 $|11\rangle \mapsto |10\rangle$

CNOT = $\begin{pmatrix} 1 & 0 & 0 & 0 \\ 0 & 1 & 0 & 0 \\ 0 & 0 & 0 & 1 \\ 0 & 0 & 1 & 0 \end{pmatrix}$

1

$|\psi\rangle$



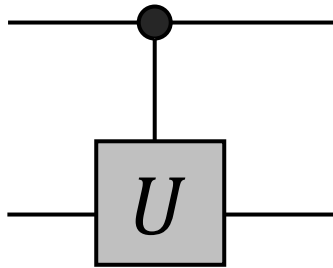
$|0\rangle$

} $|\Phi\rangle = ?$

$ \psi\rangle$	$ \Phi\rangle$
$ 0\rangle + 1\rangle$?
$ 0\rangle - 1\rangle$?

Exercise: controlled unitary

Control

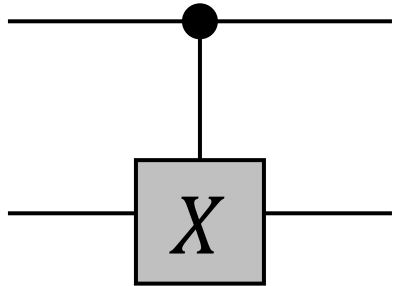


Target

$$C-U: \begin{aligned} |00\rangle &\mapsto |00\rangle \\ |01\rangle &\mapsto |00\rangle \\ |10\rangle &\mapsto |1\rangle U |0\rangle \\ |11\rangle &\mapsto |1\rangle U |1\rangle \end{aligned}$$

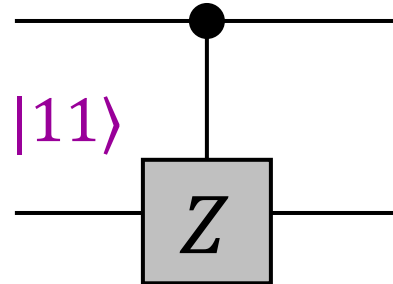
$$C-U = \begin{pmatrix} 1 & 0 & 0 & 0 \\ 0 & 1 & 0 & 0 \\ 0 & 0 & u_{00} & u_{01} \\ 0 & 0 & u_{10} & u_{11} \end{pmatrix}$$

1



$\equiv?$

2



$$|00\rangle + |11\rangle$$

$=?$

Apps of Entanglement

1. Superdense coding

How much **classical** information in n qubits?

- $2^n - 1$ complex numbers apparently needed to describe an arbitrary n -qubit state:

$$\alpha_{000}|000\rangle + \alpha_{001}|001\rangle + \alpha_{010}|010\rangle + \dots + \alpha_{111}|111\rangle$$

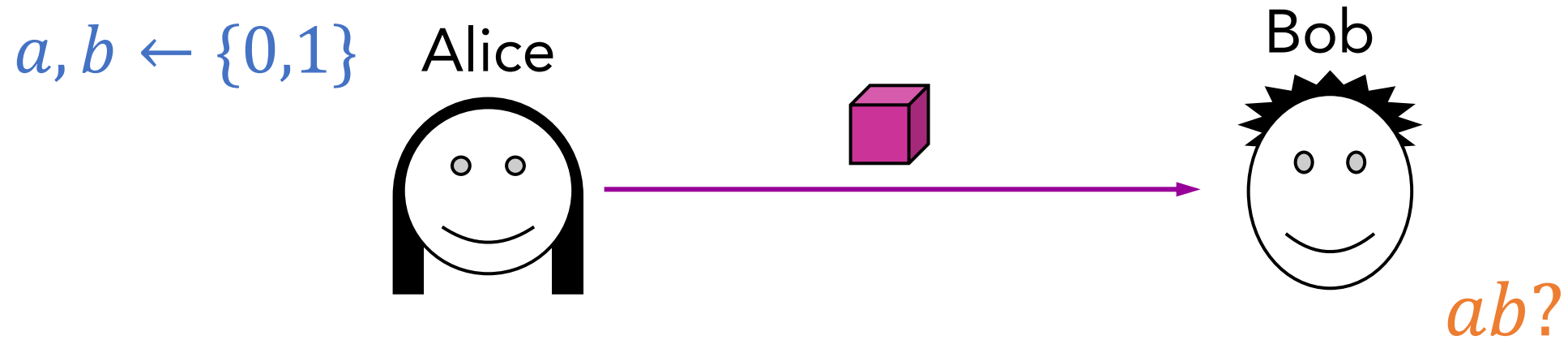
- Does this mean that an exponential amount of classical information is somehow “stored” in n qubits?

Not in an operational sense ...

Holevo's Theorem (from 1973) implies: one **cannot** convey more than n classical bits of **information** in n qubits

Superdense coding (prelude)

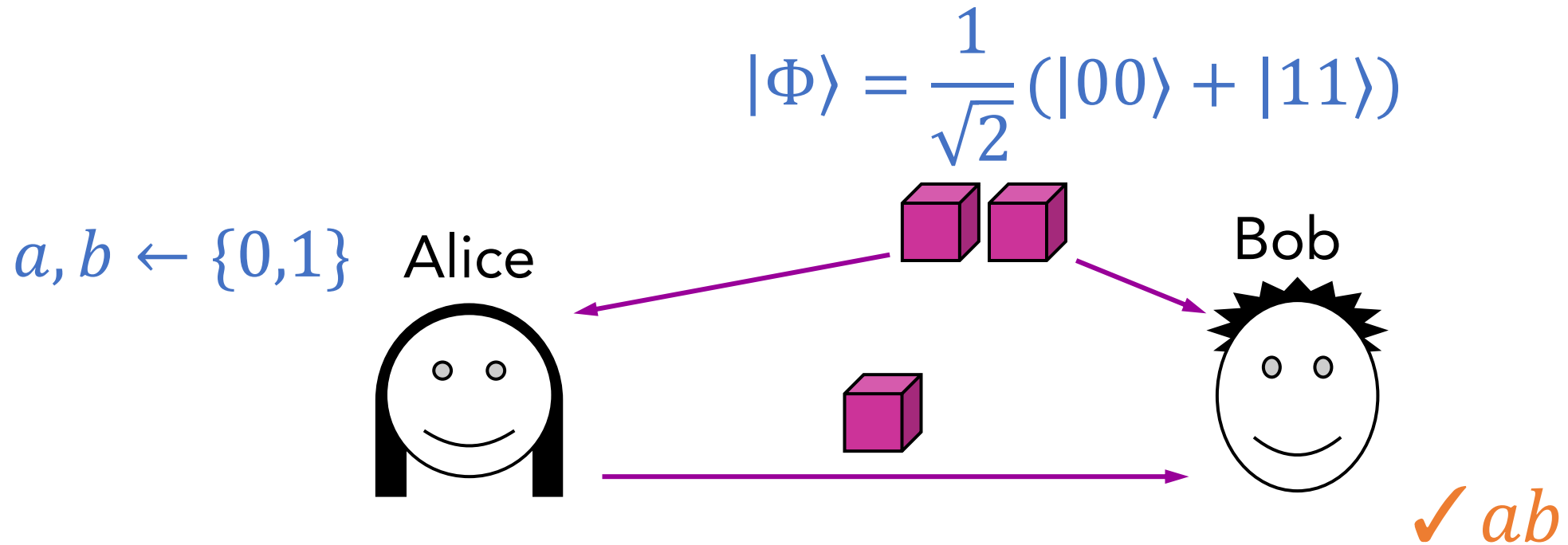
Goal: Alice wants to convey **two** classical bits to Bob sending just **one** qubit



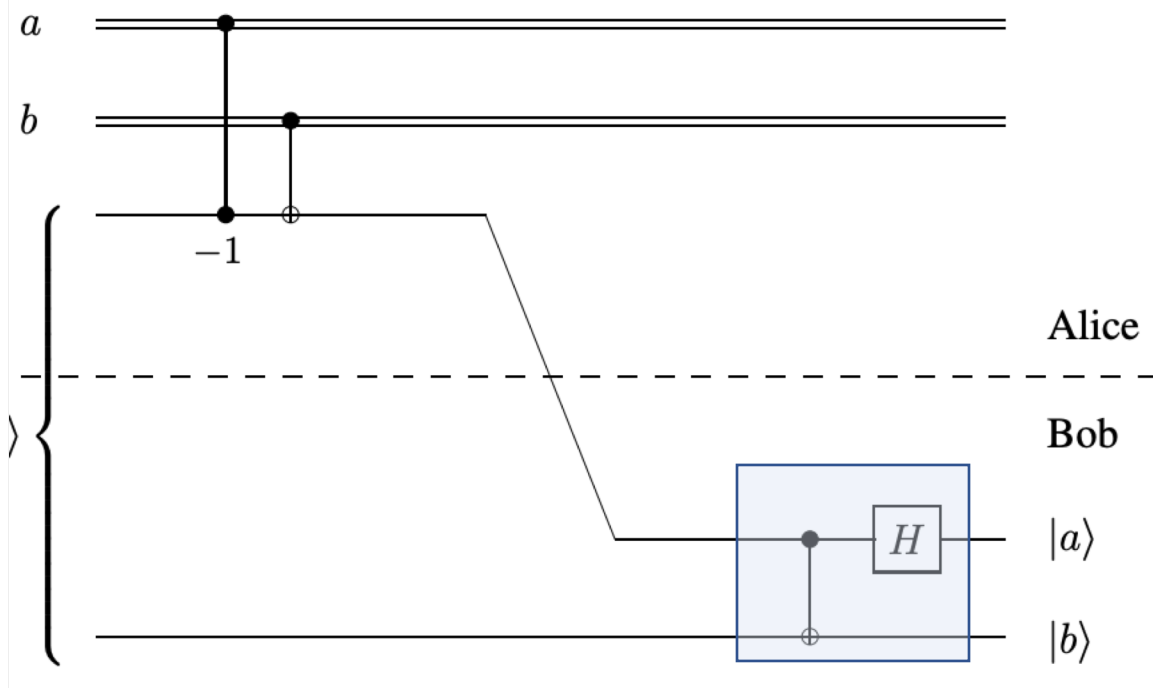
By Holevo's Theorem, this is impossible!

Superdense coding with shared EPR

Yes, if they pre-share EPR!



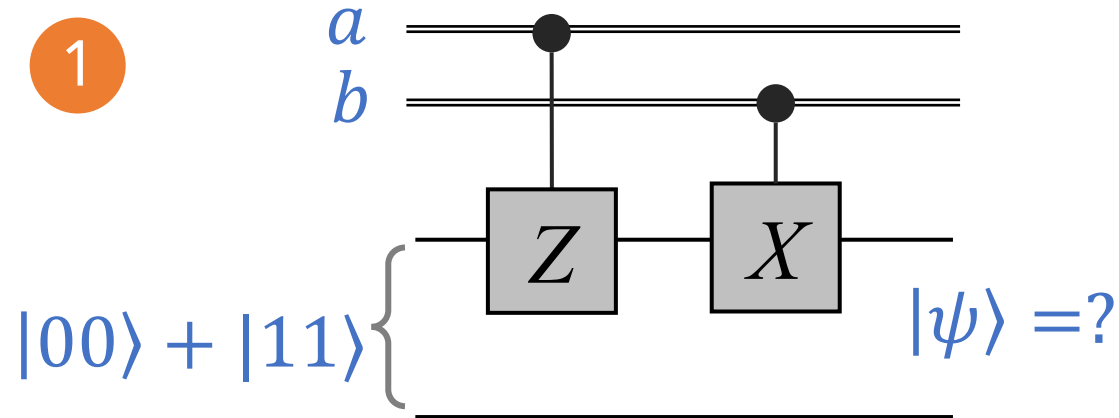
Superdense coding protocol



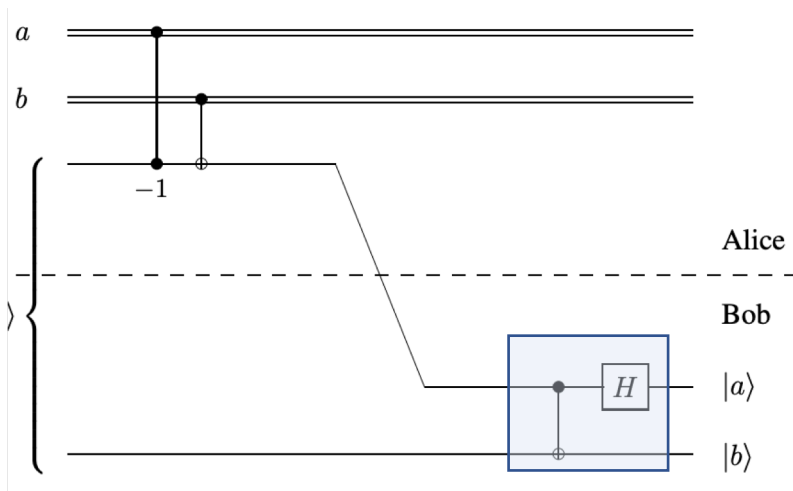
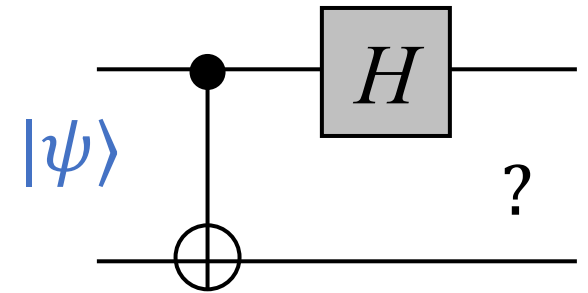
1. Bob: create $|00\rangle + |11\rangle$ and send the **first** qubit to Alice
2. Alice:
 - if $a = 1$ then apply Z to qubit
 - if $b = 1$ then apply X to qubit
 - send the qubit back to Bob
3. Bob: apply the "gadget" and measure the two qubits

Analysis

1



2



ab	$ \psi\rangle$
00	?
01	?
10	?
11	?

Bell states

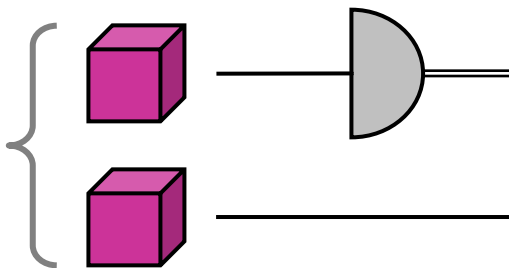
Input	Output
$ 00\rangle + 11\rangle$?
$ 01\rangle + 10\rangle$?
$ 00\rangle - 11\rangle$?
$ 01\rangle - 10\rangle$?

Apps of Entanglement

2. Quantum teleportation

Partial measurement

- Measuring the first qubit of a two-qubit system

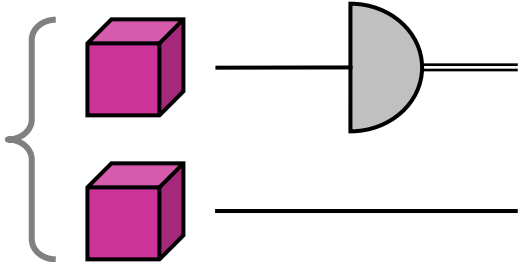
$$\alpha_{00}|00\rangle + \alpha_{01}|01\rangle + \alpha_{10}|10\rangle + \alpha_{11}|11\rangle$$


- Result

See	With probability	posterior state (renormalized!)
0	$p_0 := \alpha_{00} ^2 + \alpha_{01} ^2$	$\frac{\alpha_{00} 00\rangle + \alpha_{01} 01\rangle}{\sqrt{ \alpha_{00} ^2 + \alpha_{01} ^2}}$
1	$p_1 := \alpha_{10} ^2 + \alpha_{11} ^2$	$\frac{\alpha_{10} 10\rangle + \alpha_{11} 11\rangle}{\sqrt{ \alpha_{10} ^2 + \alpha_{11} ^2}}$

Partial measurement: Exercise

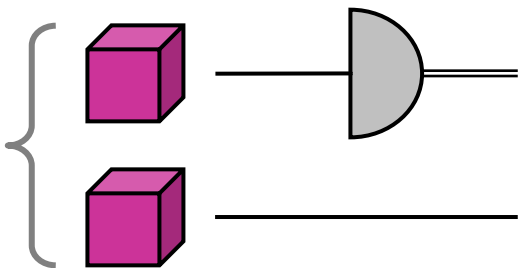
- Measuring the first qubit of a two-qubit system

$$|\psi\rangle = \frac{1}{2}|00\rangle - \frac{i}{2}|10\rangle + \frac{1}{\sqrt{2}}|11\rangle$$


See	With probability	posterior state (renormalized!)
0		
1		

Partial measurement: Exercise

- Measuring the first qubit of a two-qubit system

$$|\psi\rangle = \frac{1}{2}|00\rangle - \frac{i}{2}|10\rangle + \frac{1}{\sqrt{2}}|11\rangle$$


- A trick

See	With probability	posterior state (renormalized!)
0		
1		

Transmitting qubits by classical bits

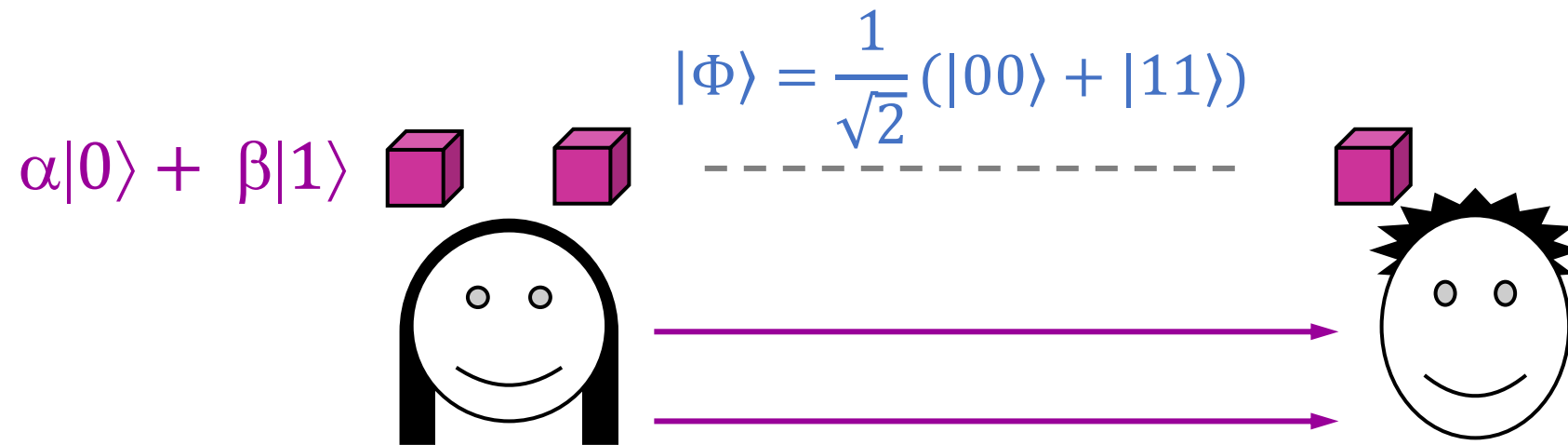
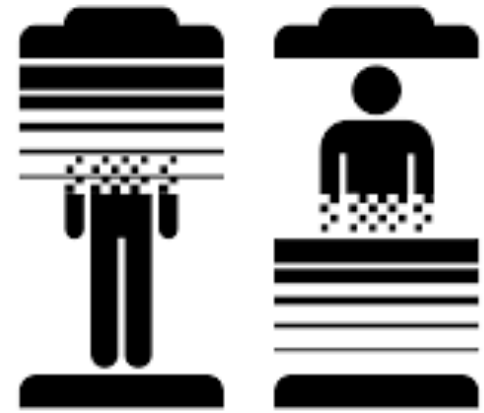
Goal: Alice conveys a **qubit** to Bob by sending just **classical** bits



- If Alice **knows** $\alpha, \beta \in \mathbb{C}$, requires **infinitely many bits** for perfect precision
- If Alice **doesn't know** α or β , she can at best acquire **one bit** by measurement

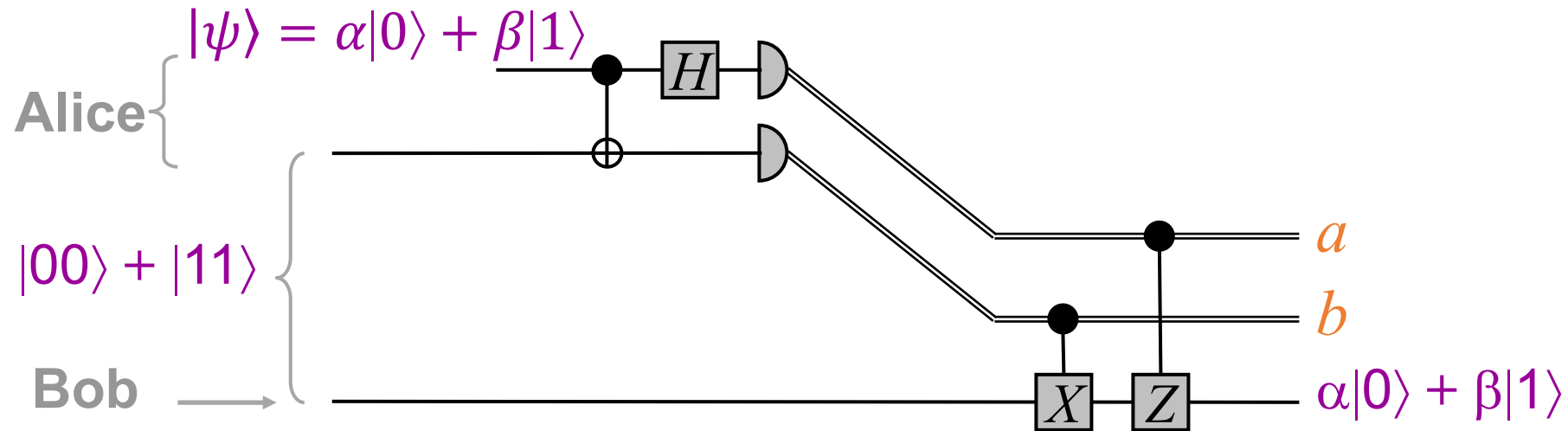
Teleportation

Theorem. **two** classical bit enough if pre-share **EPR**



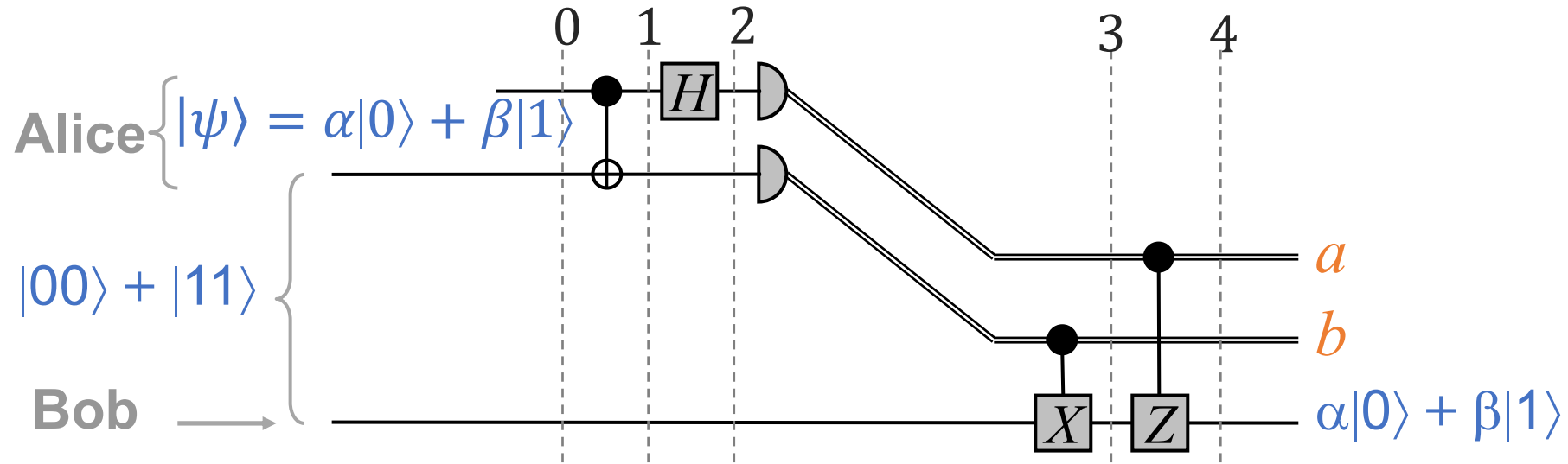
Teleportation: protocol

Theorem. **two** classical bit enough if pre-share **EPR**



- Does Alice still hold $|\psi\rangle$ at the end?
- Communicating faster than the speed of light?

Teleportation: analysis



Questions?

- Use zoom chat and campuswire DM/chatroom to mingle and identify potential group members
- Ask me if you want a Zoom breakout room

